

# Conditioning analysis of a continuous time subspace-based model identification algorithm

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## Abstract

We present in this paper a study concerning the conditioning analysis of a continuous-time deterministic subspace-based model identification algorithm. We show that the conditioning number of the associated extended observability matrices depends on an exponential way from both: the estimated order of the system and the dimension of the system output vector.

## 1 Introduction

As far as linear time-invariant multivariable systems are concerned, subspace-based model identification algorithms constitute a broad family of identification methods mainly characterized by the use of geometric information (see for instance [3], [4], and [5]). The subspace-based methods compute the estimated parameters, i.e., a state space realization of the system, from an approximation of the observability subspace of the concerned system. This approximation is obtained from a discrete-time set of input-output measurements. Continuous-time data is filtered and sampled in order to obtain discrete-time information (see [1]). We present in this paper an algorithm to identify a continuous linear time-invariant model from a given set of discrete-time input-output measurements. Our algorithm is based on the method proposed in [3], and depends essentially on the pseudo-inverse of the so-called extended observability matrices. We present then the analysis of the conditioning number of these observability matrices, which is the main purpose of this paper.

In Section 2 we present our proposed algorithm, while in Section 3 we discuss its numerical properties.

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## 2 The CN4SID algorithm

Consider the continuous linear time-invariant system  $(A, B, C, D)$  described by:

$$\begin{cases} \mathbf{p}x(t) = Ax(t) + Bu(t), \\ y(t) = Cx(t) + Du(t), \end{cases} \quad (2.1)$$

where:  $\mathbf{p}$  denotes the differential operator, i.e.  $\mathbf{p} = d/dt$ ;  $x(\cdot) \in \mathbb{R}^n$  denotes the state;  $u(\cdot) \in \mathbb{R}^m$  denotes the input, and  $y(\cdot) \in \mathbb{R}^p$  denotes the output.  $A, B, C$  and  $D$  are linear maps represented by real constant matrices. It is assumed that  $(C, A)$  is observable.

Consider the scalar causal stable operator:

$$\lambda = \frac{1}{1 + \mathbf{p}\tau}, \quad (2.2)$$

where  $\tau$  is a scalar such that  $\tau > 0$ . Given a  $\mathbb{R}^p$  valued signal  $y(t)$ , we define the  $i^{\text{th}}$ -filtered signal  $[\lambda^i y](t)$  as follows:

$$[\lambda^i y](t) = \begin{cases} \lambda [\lambda^{i-1} y](t), & \text{for } i > 1, \\ y(t), & \text{when } i = 0. \end{cases}$$

Applying this filtering action to the system  $(A, B, C, D)$  we obtain the modified system:

$$\begin{cases} x(t) = A_\lambda [\lambda x](t) + B_\lambda [\lambda u](t), \\ y(t) = Cx(t) + Du(t), \end{cases} \quad (2.3)$$

where  $A_\lambda = I + \tau A$  and  $B_\lambda = \tau B$ . Since  $(A, B)$  and  $(A_\lambda, B_\lambda)$  are linked through a bijective map, the obtention of estimates for  $A$  and  $B$  results in the obtention of estimates for  $A$  and  $B$ .

Suppose constant sampling and a sampling time sequence  $\{t_k\}_{k=0}^{N-1}$  be given, where  $t_k = t_0 + kh$  (with  $h$  denoting the sampling period). Let input-output measurements  $\{u_k\}_{k=0}^{N-1}$  and  $\{y_k\}_{k=0}^{N-1}$  on system (2.3) be given. The deterministic **C**ontinuous **S**ubspace **I**dentification (**CSId**) problem is then defined as follows:

**Definition 2.1. CSId problem:** Consider the filtered model (2.3) and the input-output measurements  $\{u_k\}_{k=0}^{N-1}$  and  $\{y_k\}_{k=0}^{N-1}$ , estimate then the unknown matrices  $A, B, C$ , and  $D$ .

We can then build the following matrices:

$$Y := \begin{bmatrix} [\lambda^{i-1} y]_0 & [\lambda^{i-1} y]_1 & \cdots & [\lambda^{i-1} y]_{N-1} \\ [\lambda^{i-2} y]_0 & [\lambda^{i-2} y]_1 & \cdots & [\lambda^{i-2} y]_{N-1} \\ \vdots & \vdots & \ddots & \vdots \\ y_0 & y_1 & \cdots & y_{N-1} \end{bmatrix} \quad (2.4)$$

and:

$$U := \begin{bmatrix} [\lambda^{i-1}u]_0 & [\lambda^{i-1}u]_1 & \cdots & [\lambda^{i-1}u]_{N-1} \\ [\lambda^{i-2}u]_0 & [\lambda^{i-2}u]_1 & \cdots & [\lambda^{i-2}u]_{N-1} \\ \vdots & \vdots & \ddots & \vdots \\ u_0 & u_1 & \cdots & u_{N-1} \end{bmatrix}. \quad (2.5)$$

In what follows we present the **CN4SID** algorithm, which is based on: the **N4SID** algorithm [3], and the space transformation discussed in [1].

**CN4SID Algorithm:** Consider a continuous linear-time invariant system described by (2.1) and the corresponding filtered model (2.3). Let the input-output sequences  $\{u_k\}_{k=0}^{N-1}$  and  $\{y_k\}_{k=0}^{N-1}$  be given. Let  $i$  be an the estimated order of the system such that  $N \gg i > n$ , where  $n$  denotes the real system. It is then possible to find matrices  $\widehat{A}$ ,  $\widehat{B}$ ,  $\widehat{C}$ , and  $\widehat{D}$  (i.e., the estimates of  $A$ ,  $B$ ,  $C$ , and  $D$ , respectively), according to the procedure:

1. Build matrices  $Y$  and  $U$  as defined in (2.4) and (2.5), respectively.
2. Build matrices  $[\lambda^i Y]$  and  $[\lambda^i U]$ .
3. Build matrices  $[\lambda^i W]$  and  $[\lambda^{i+1} W]$  defined as follows:

$$[\lambda^i W] := \begin{bmatrix} [\lambda^i U] \\ [\lambda^i Y] \end{bmatrix}, \text{ and } [\lambda^{i+1} W] := \begin{bmatrix} [\lambda^{i+1} U] \\ [\lambda^{i+1} Y] \end{bmatrix}.$$

4. Compute the projections:

$$Y / \begin{bmatrix} [\lambda^i W] \\ [\lambda^i U] \end{bmatrix} = [ L_w \quad L_u ] \begin{bmatrix} [\lambda^i W] \\ [\lambda^i U] \end{bmatrix}$$

and:

$$[\lambda Y] / \begin{bmatrix} [\lambda^{i+1} W] \\ [\lambda^{i+1} U] \end{bmatrix} = [ L_{\bar{w}} \quad L_{\bar{u}} ] \begin{bmatrix} [\lambda^{i+1} W] \\ [\lambda^{i+1} U] \end{bmatrix}.$$

5. From the singular value decomposition of  $L_w$ :

$$\Lambda_i X_i = L_w [\lambda^i W] = [ U_N \quad U_0 ] \begin{bmatrix} S_n & 0 \\ 0 & S_0 \end{bmatrix} \begin{bmatrix} V'_n \\ V'_0 \end{bmatrix} [\lambda^i W]$$

compute the  $i^{th}$ - $\Lambda$ -extended observability matrices:

$$\Lambda_i = U_n S_n^{1/2}, \quad (2.6)$$

where  $S_n^{1/2}$  stands for the Cholesky factor of  $S_n$ .

6. From extended observability matrices  $\Lambda_i$  compute the state sequences:

$$X_i = \Lambda_i^\dagger L_w [\lambda^i W]$$

and:

$$[\lambda X_i] = \Lambda_i^\dagger L_w [\lambda^{i+1} W].$$

$\Lambda_i^\dagger$  stands for the pseudo-inverse of  $\Lambda_i$ .

7. Solve the following optimization problem:

$$\min_k \left\| \begin{bmatrix} X_i \\ [\lambda^i y] \end{bmatrix} - K \begin{bmatrix} [\lambda X_i] \\ [\lambda^i u] \end{bmatrix} \right\|_F^2,$$

where:

$$K := \begin{bmatrix} A_\lambda & B_\lambda \\ \widehat{C} & \widehat{D} \end{bmatrix}.$$

and  $\|\alpha\|_F$  stands for the Frobenius norm of  $\alpha$ .

8. Finally:

$$\widehat{A} = \frac{A_\lambda - I}{\tau}$$

and:

$$\widehat{B} = \frac{B_\lambda}{\tau}.$$

In what follows we discuss the numerical properties of the **CN4SID** algorithm.

### 3 Numerical properties

Since the **CN4SID** algorithm depends essentially on the pseudo-inverse of the  $\Lambda$ -extended observability matrices  $\Lambda_i$  given by (2.6), in this section we discuss the conditioning of these matrices.

We first present some preliminary results concerning the nature of  $\Lambda_i$ . Let us define the  $i^{\text{th}}$ -extended observability matrix  $\Gamma_i$  as follows:

$$\Gamma_i := \begin{bmatrix} C \\ CA \\ CA^2 \\ \vdots \\ CA^{i-1} \end{bmatrix}. \quad (3.7)$$

Then:

**Lemma 3.1.** *Let  $\Gamma_i$  be as defined in (3.7), then:*

$$\Lambda_i = T\Gamma_i,$$

where:

$$T = \begin{bmatrix} I & 0 & 0 & \cdots & 0 \\ I & \tau I & 0 & \cdots & 0 \\ I & 2\tau I & \tau^2 I & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ I & \binom{i-1}{1}\tau I & \binom{i-1}{2}\tau^2 I & \cdots & \tau^{i-1} I \end{bmatrix} \in \mathbb{R}^{ip \times ip}, \quad (3.8)$$

with:

$$\binom{v}{z} := \frac{v!}{(v-z)!z!}.$$

*Proof.* Because of the definition of  $\Lambda_i$  we have that:

$$\Lambda_i = \begin{bmatrix} C \\ CA_\lambda \\ \vdots \\ CA_\lambda^{i-1} \end{bmatrix}, \quad (3.9)$$

where:

$$A_\lambda = I + \tau A. \quad (3.10)$$

Thus:

$$\Lambda_i = \begin{bmatrix} C \\ C + \tau CA \\ C + 2\tau CA + \tau^2 CA^2 \\ \vdots \\ C + \binom{i-1}{1}\tau CA + \binom{i-1}{2}\tau^2 CA^2 + \cdots + \tau^{i-1} CA^{i-1} \end{bmatrix} \quad (3.11)$$

which implies:

$$\Lambda_i = T\Gamma_i. \quad (3.12)$$

□

Thus, the numerical conditioning of  $\Lambda_i$ , in terms of the Frobenius norm, is given by:

$$k_F(\Lambda_i) = \|\Lambda_i\|_F \|\Lambda_i^{-1}\|_F,$$

and from Lemma 3.1:

$$k_F(\Lambda_i) = \|T\Gamma_i\|_F \|(T\Gamma_i)^{-1}\|_F.$$

Consequently:

$$k_F(\Lambda_i) \leq \|T\|_F \|T^{-1}\|_F \|\Gamma_i\|_F \|\Gamma_i^{-1}\|_F,$$

i.e.:

$$k_F(\Lambda_i) \leq k_F(T) k_F(\Gamma_i). \quad (3.13)$$

**Remark 3.1.** *Because of its dependence on the input-output measurements, it is not possible to have a priori knowledge of  $k_F(\Gamma_i)$ . However, it is suitable to have a small value for this conditioning number. As far as  $k_F(T)$  is concerned, it is possible to have a priori knowledge of its value. For the above reason, we focus our analysis on the study of the conditioning number of  $T$ .*

**Lemma 3.2.** *Let  $T$  be as defined in (3.8), then:*

$$\|T\|_F^2 = p \sum_{k=1}^i \sum_{j=1}^k \binom{k-1}{j-1}^2 \tau^{2(j-1)}.$$

*Proof.* Let matrix  $T$  be as defined in (3.8). Then, the block matrix  $T_{kj} \in \mathbb{R}^{p \times p}$  can be written as follows:

$$T_{kj} = \binom{k-1}{j-1} \tau^{j-1} I_{p \times p} \quad (3.14)$$

and its corresponding transpose block matrix is given by:

$$T'_{jl} = \binom{l-1}{j-1} \tau^{j-1} I_{p \times p} \quad (3.15)$$

and consequently the  $kl$ -block matrix of  $TT'$ , i.e.  $(TT')_{kl} \in \mathbb{R}^{p \times p}$ , is given by:

$$\begin{aligned} (TT')_{kl} &= \sum_{j=1}^i \binom{k-1}{j-1} \tau^{j-1} I_{p \times p} \binom{l-1}{j-1} \tau^{j-1} I_{p \times p} \\ &= \sum_{j=1}^i \binom{k-1}{j-1} \binom{l-1}{j-1} \tau^{2(j-1)} I_{p \times p}. \end{aligned} \quad (3.16)$$

Moreover:

$$\begin{aligned} \text{tr}(TT') &= \sum_{k=l=1}^i \sum_{j=1}^i \binom{k-1}{j-1} \binom{l-1}{j-1} \tau^{2(j-1)} p \\ &= \sum_{k=1}^i \sum_{j=1}^k \binom{k-1}{j-1}^2 \tau^{2(j-1)} p, \end{aligned} \quad (3.17)$$

and because of the definition of the Frobenius norm:

$$\|T\|_F^2 = \text{tr}(TT') \quad (3.18)$$

$$= p \sum_{k=1}^i \sum_{j=1}^k \binom{k-1}{j-1}^2 \tau^{2(j-1)} \quad (3.19)$$

□

**Remark 3.2.** From the previous lemma we have that  $\|T\|_F^2$  is a  $2(j-1)$  polynomial function of  $\tau$ . It is usual that  $0 < \tau < 1$ . Thus,  $\|T\|_F^2$  usually has a small value.

As far as the Frobenius norm of  $T^{-1}$  is concerned, we have the following result, which gives the Frobenius norm of  $\text{adj}(T)$ , i.e., the adjunct of  $T$ :

**Lemma 3.3.** Let  $T$  be as defined in (3.8), then:

$$\|\text{adj}(T)\|_F^2 = p\tau^{n(n-1)} \sum_{k=1}^i \sum_{j=1}^k \binom{k-1}{j-1}^2 \tau^{-2(k-1)},$$

where  $p > 0$  denotes the dimension of the output vector.

*Proof.* Let matrix  $T$  be as defined in (3.8). Then, the adjunct of  $T$  can be written in terms of its  $kj$ -block matrix, i.e.  $(\text{adj}T)_{kj} \in \mathbb{R}^{p \times p}$ , as follows:

$$(\text{adj}T)_{kj} = (-1)^{k+j} \binom{k-1}{j-1} \tau^{\frac{i(i-1)}{2} - (k-1)} I_{p \times p} \quad (3.20)$$

and the corresponding block matrices of the transpose of  $\text{adj}(T)$  are given by:

$$(\text{adj}T)'_{jl} = (-1)^{l+j} \binom{l-1}{j-1} \tau^{\frac{i(i-1)}{2} - (l-1)} I_{p \times p}. \quad (3.21)$$

Then, the block matrices of the matrix product  $\text{adj}(T) \text{adj}(T)'$  can be written as:

$$\begin{aligned} (\text{adj}T (\text{adj}T'))_{kl} &= \sum_{j=1}^i (-1)^{k+j} \binom{k-1}{j-1} \tau^{\frac{i(i-1)}{2} - (k-1)} I_{p \times p} \\ &\quad \cdot (-1)^{l+j} \binom{l-1}{j-1} \tau^{\frac{i(i-1)}{2} - (l-1)} I_{p \times p} \\ &= \sum_{j=1}^i (-1)^{k+l+2j} \binom{k-1}{j-1} \binom{l-1}{j-1} \tau^{i(i-1)-2(k-1)} I_{p \times p}. \end{aligned} \quad (3.22)$$

As far as the trace of the matrix product  $\text{adj}(T) \text{adj}(T)'$  is concerned we have:

$$\begin{aligned} \text{tr}(\text{adj}(T) \text{adj}(T)') &= \sum_{k=1}^i \sum_{j=1}^i (-1)^{k+l+2j} \binom{k-1}{j-1} \\ &\quad \cdot \binom{l-1}{j-1} \tau^{\frac{i(i-1)}{2} - (k-1)} \tau^{\frac{i(i-1)}{2} - (l-1)} p \\ &= \sum_{k=1}^i \sum_{j=1}^i \binom{k-1}{j-1}^2 \tau^{i(i-1)-2(k-1)} p. \end{aligned} \quad (3.23)$$

Consequently:

$$\begin{aligned}
\|adj(T)\|_F^2 & : = tr(adj(T) adj(T)') \\
& = p\tau^{i(i-1)} \sum_{k=1}^i \sum_{j=1}^k \binom{k-1}{j-1}^2 \tau^{-2(k-1)}.
\end{aligned} \tag{3.24}$$

□

We can at this level present our main result. Combining Lemma 3.2 and Lemma 3.3 we have:

**Theorem 3.1.** *Let  $T$  be as defined in (3.8), then:*

$$k_F(T) = p\tau^{\frac{i(i-1)(1-p)}{2}} \sqrt{\left( \sum_{k=1}^i \sum_{j=1}^k \binom{k-1}{j-1}^2 \tau^{2(j-1)} \right) \left( \sum_{k=1}^i \sum_{j=1}^k \binom{k-1}{j-1}^2 \tau^{-2(j-1)} \right)}.$$

*Proof.* By definition:

$$k_F(T) := \|T\|_F \|T^{-1}\|_F, \tag{3.25}$$

i.e.:

$$\begin{aligned}
k_F(T) & = \sqrt{tr(TT') tr(T^{-1}(T^{-1})')} \\
& = \sqrt{tr(TT') \frac{tr(adj(T) adj(T)')}{\det(TT')}}.
\end{aligned} \tag{3.26}$$

Now, taking into account Lemma 3.2, Lemma 3.3, and since  $\det(T) = \tau^{\frac{ip(i-1)}{2}}$ , we have:

$$\begin{aligned}
& k_F(T) \\
& = \sqrt{\left( p \sum_{k=1}^i \sum_{j=1}^k \binom{k-1}{j-1}^2 \tau^{2(j-1)} \right) \left( p\tau^{i(i-1)} \sum_{k=1}^i \sum_{j=1}^k \binom{k-1}{j-1}^2 \tau^{-2(k-1)} \right) \tau^{ip(i-1)}} \\
& = p\tau^{\frac{i(i-1)(1-p)}{2}} \sqrt{\left( \sum_{k=1}^i \sum_{j=1}^k \binom{k-1}{j-1}^2 \tau^{2(j-1)} \right) \left( \sum_{k=1}^i \sum_{j=1}^k \binom{k-1}{j-1}^2 \tau^{-2(k-1)} \right)}.
\end{aligned} \tag{3.27}$$

□

**Remark 3.3.** *As is established in Theorem 3.1, the conditioning number of  $T$  increases (and so the conditioning number of  $\Lambda_i$ ) in an exponential way with both the estimated order of the system,  $i$ , and the number of outputs,  $p$ .*



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